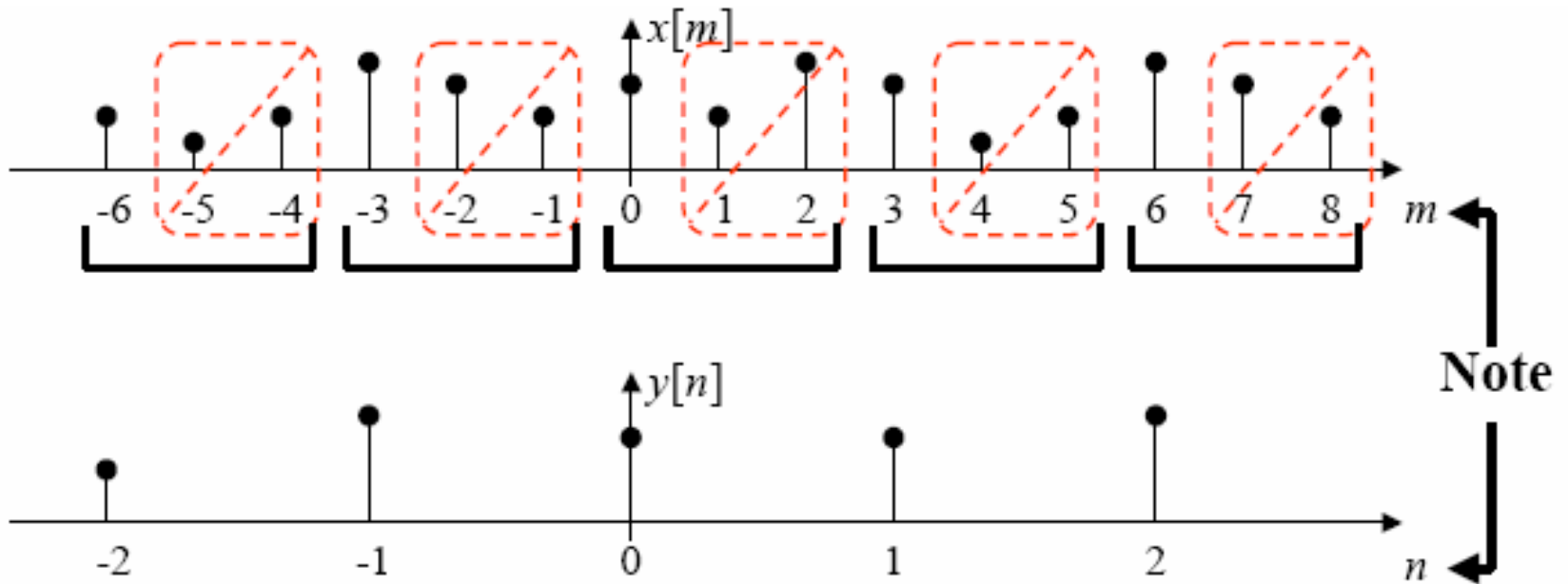


Decimation by M



Time domain representation





Decimator by M

- Z-transform analysis of decimator (continued)

$$U(z) \quad \rightarrow \quad \boxed{\downarrow M} \quad \rightarrow \quad \frac{1}{M} \sum_{i=0}^{M-1} U(z^{1/M} e^{-j2\pi i/M})$$

- Note that $U(e^{j\omega})$ is periodic with period 2π , while $U(e^{j\omega/M})$ is periodic with period $2M\pi$
- the summation with $i=0\dots M-1$ restores the periodicity with period 2π

- Example:

$$u[k] = \alpha^k, k \geq 0$$

$$\Rightarrow U(z) = \frac{1}{1 - \alpha z^{-1}}$$

$$\Rightarrow Y(z) = \frac{1}{M} \sum_{i=0}^{M-1} (\dots) = \dots = \frac{1}{1 - \alpha^M z^{-1}}$$

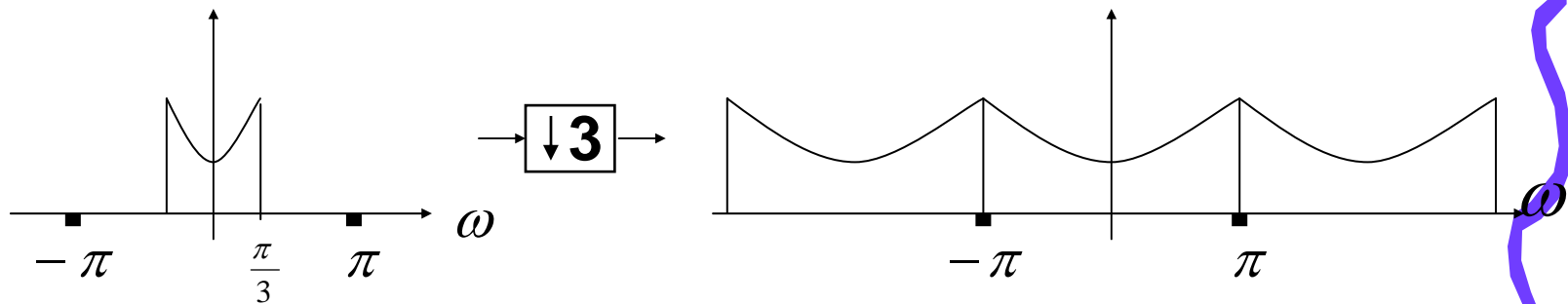
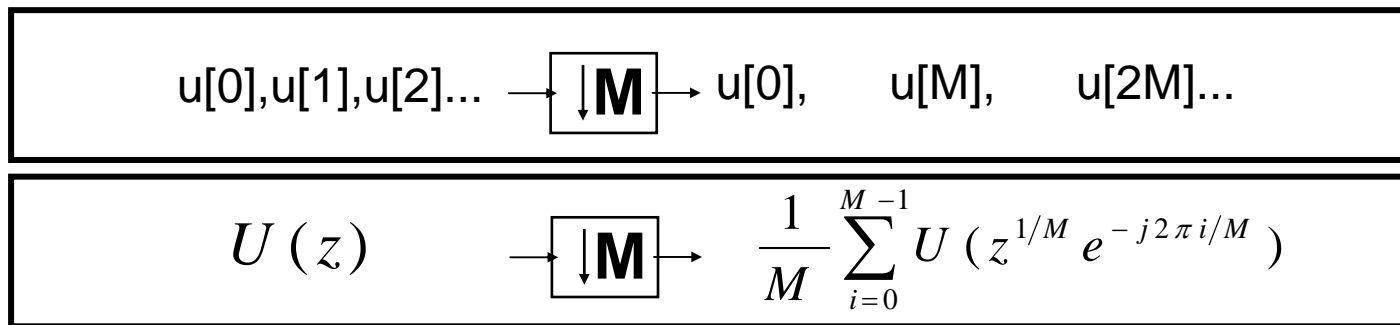
$$\Rightarrow y[k] = (\alpha^M)^k, k \geq 0$$





Decimation by M

- Z-transform (frequency domain) analysis of decimator

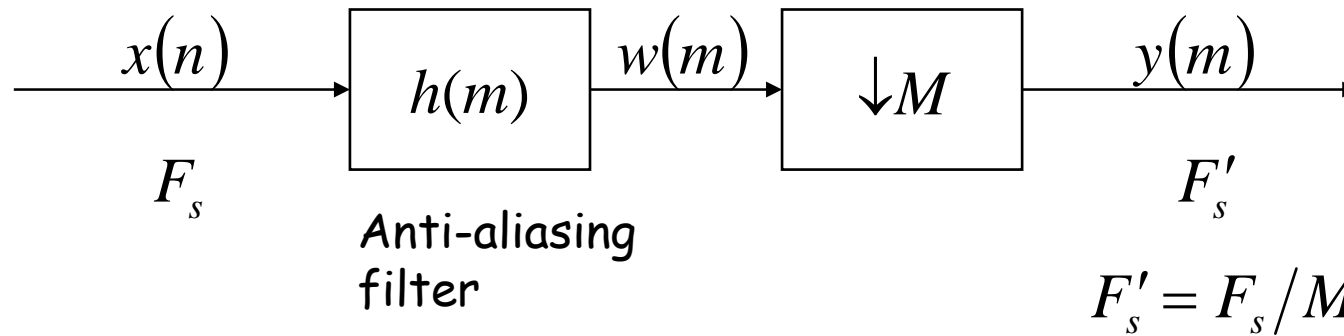


decimation introduces **ALIASING** if input signal occupies frequency band larger than π / M , for $-\pi \leq \omega \leq \pi$
 hence decimation mostly preceded by **anti-aliasing (decimation) filter**



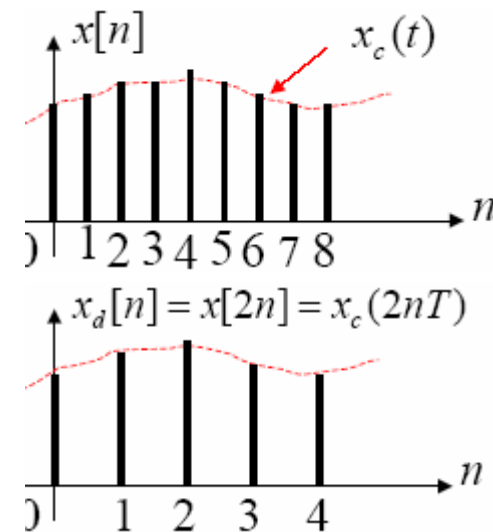


Decimation by M



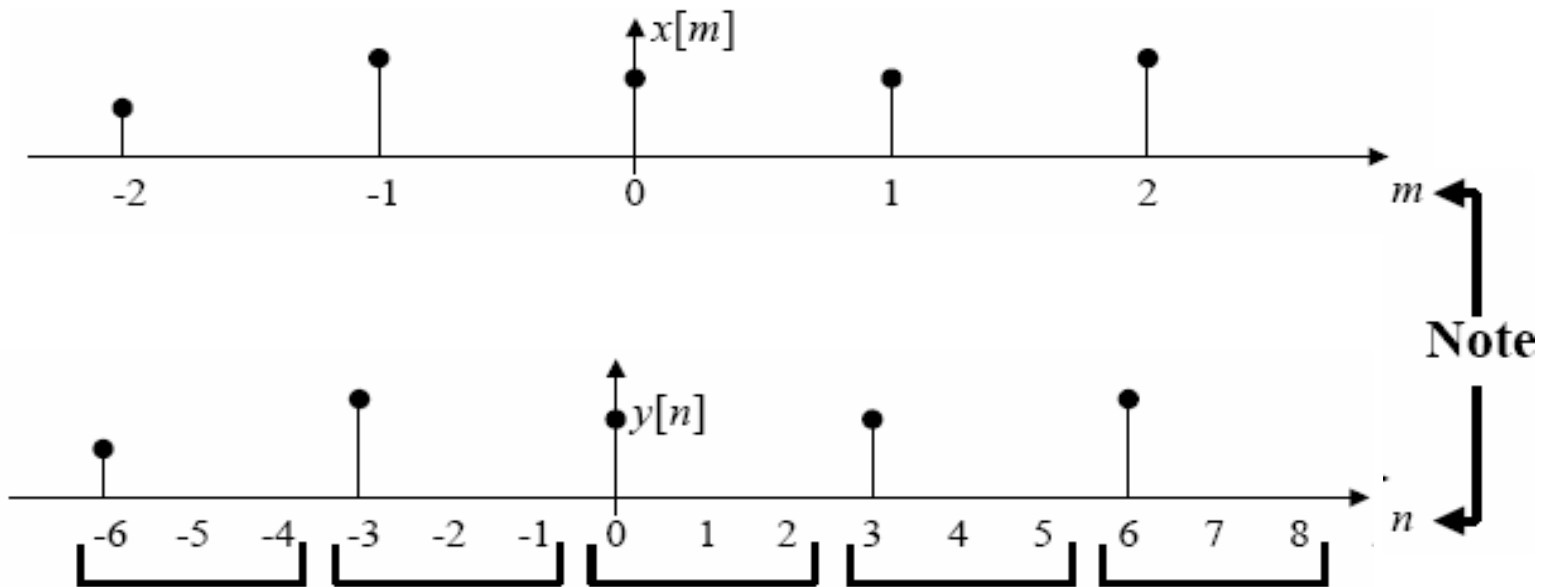
$$\frac{T'}{T} = M = \frac{F_s}{F'_s}$$

$$H_I(\omega) = \begin{cases} 1, & |\omega| < \frac{\pi}{M} \\ 0, & \text{otherwise} \end{cases}$$





Interpolation by L



Time domain representation



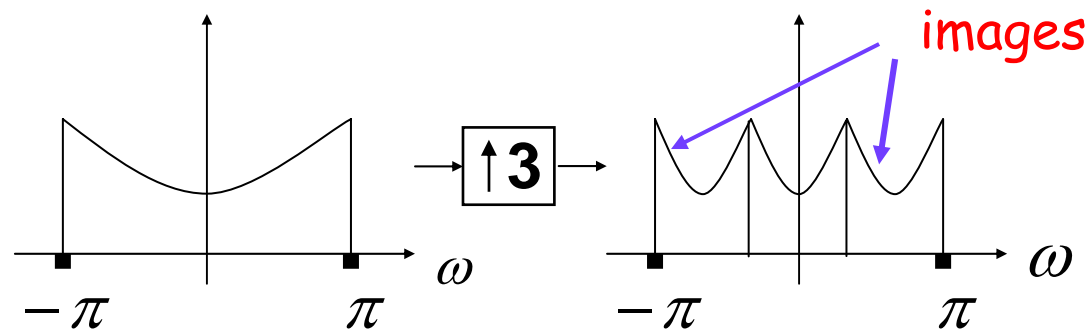


Interpolation by L

- Z-transform (frequency domain) analysis of expander

$$u[0], u[1], u[2], \dots \rightarrow \boxed{\uparrow L} \rightarrow u[0], 0, \dots, 0, u[1], 0, \dots, 0, u[2], \dots$$

$$U(z) \rightarrow \boxed{\uparrow L} \rightarrow U(z^N)$$

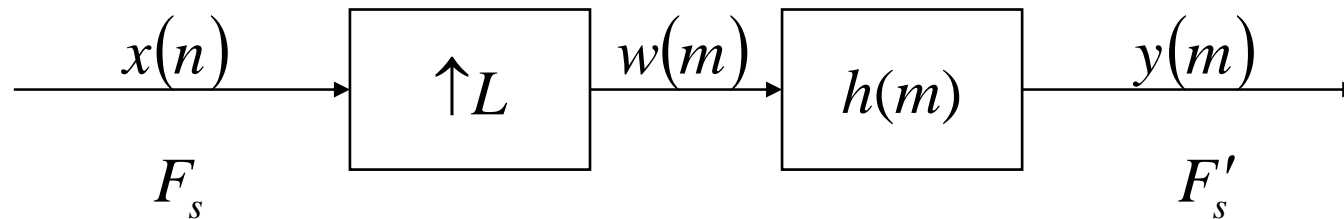


'expansion in time domain ~ compression in frequency domain'
 expander mostly followed by **interpolation filter** to remove all **images**





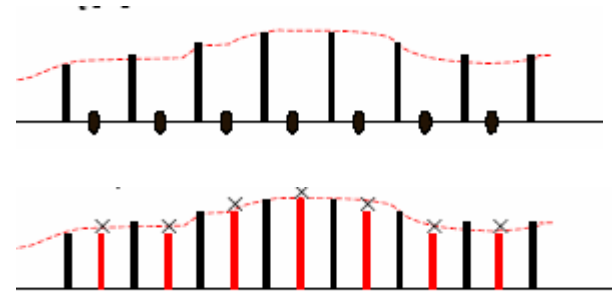
Interpolation by L



$$F'_s = LF_s$$

$$\frac{T'}{T} = \frac{1}{L} = \frac{F_s}{F'_s}$$

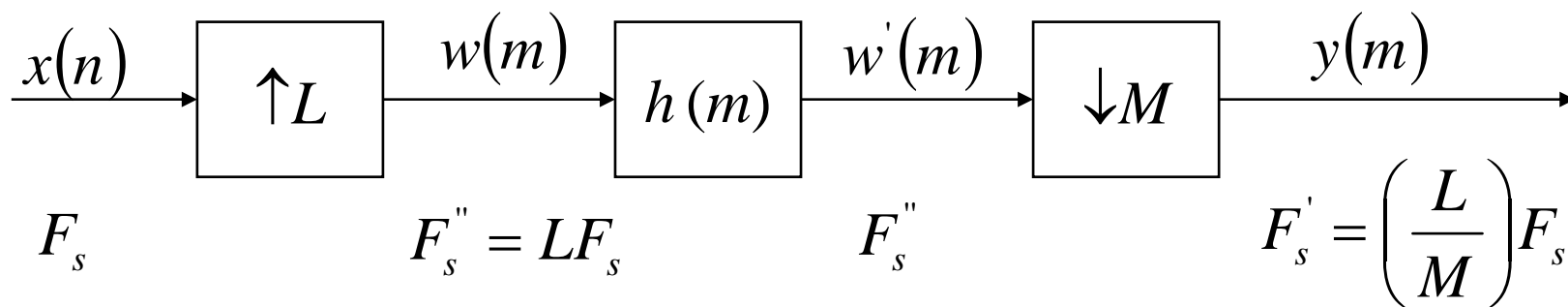
$$H_I(\omega') = \begin{cases} G, & |\omega'| < \frac{\pi}{L} \\ 0, & \text{otherwise} \end{cases}$$





Conversion by a Rational Factor M/L

- A more efficiency implementation

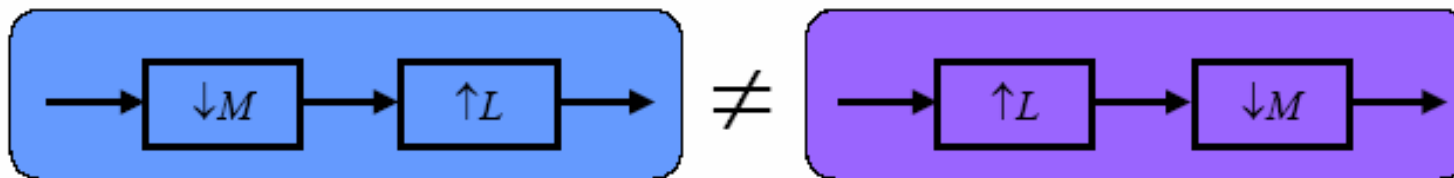
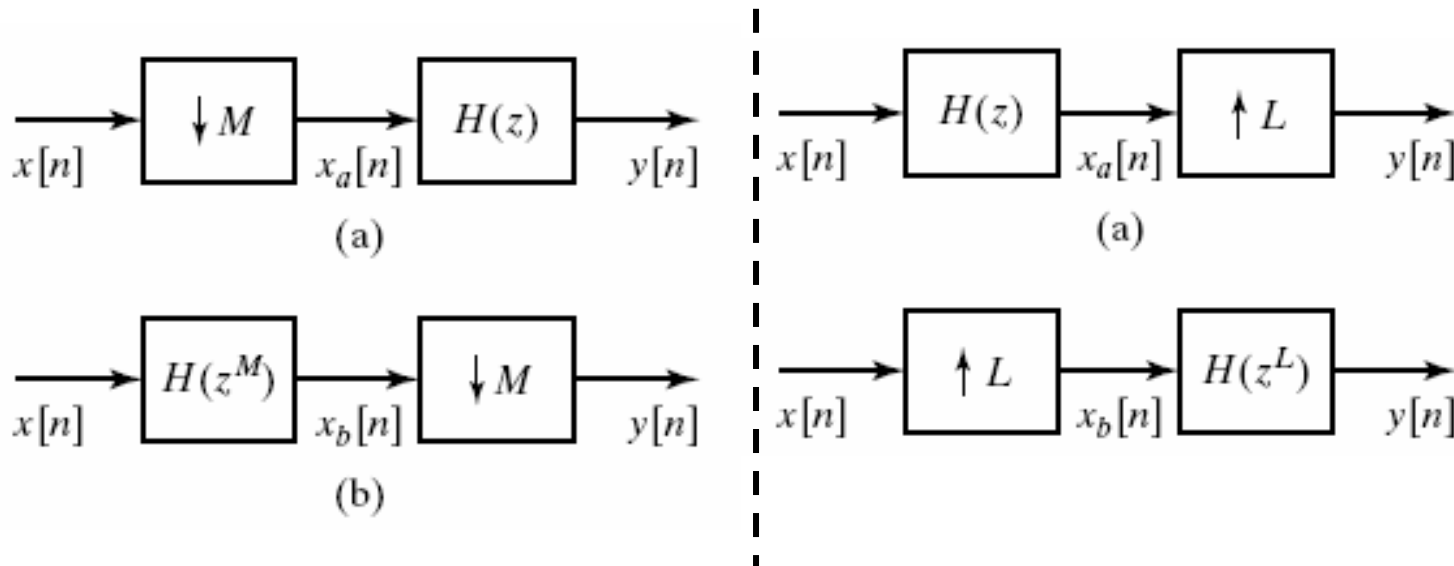


$$H_I(\omega'') = \begin{cases} L, & |\omega''| \leq \min\left\{\frac{\pi}{L}, \frac{\pi}{M}\right\} \\ 0, & \text{otherwise} \end{cases}$$





Interchange of Filtering





Proof

(b) $\implies X_b(e^{j\omega}) = H(e^{j\omega M})X(e^{j\omega})$. Recall that $Y(e^{j\omega}) = \frac{1}{M} \sum_{i=0}^{M-1} X_b(e^{j(\omega/M - 2\pi i/M)})$

$$\begin{aligned} \text{Then } Y(e^{j\omega}) &= \frac{1}{M} \sum_{i=0}^{M-1} H(e^{j(\omega - 2\pi i)}) X(e^{j(\omega/M - 2\pi i/M)}) \\ &= \frac{1}{M} \sum_{i=0}^{M-1} H(e^{j\omega}) X(e^{j(\omega/M - 2\pi i/M)}) \\ &= H(e^{j\omega}) X_a(e^{j\omega}) \end{aligned}$$

Similarly, (a) $\implies Y(e^{j\omega}) = X_a(e^{j\omega L}) = X(e^{j\omega L})H(e^{j\omega L}) = H(e^{j\omega L})X_b(e^{j\omega})$

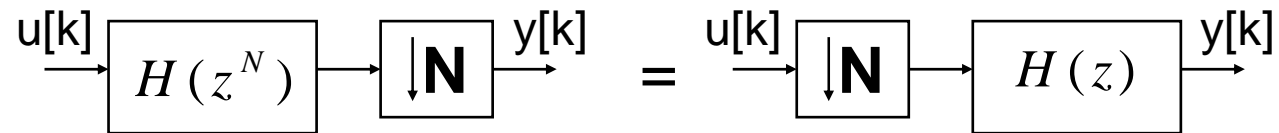
Since $X_b(e^{j\omega}) = X(e^{j\omega L})$





Identities (I)

- Nobel identity (I): (only for rational functions)



Example : N=2

$h[0], h[1], 0, 0, 0, \dots$

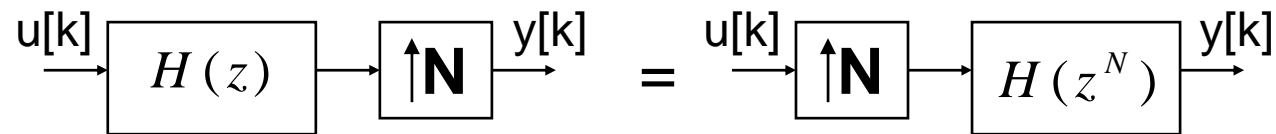
$$\begin{aligned}
 \begin{bmatrix} y[0] \\ y[1] \\ y[2] \end{bmatrix} &= \underbrace{\begin{bmatrix} 1 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & 0 \end{bmatrix}}_{\text{2-fold downsampling}} \cdot \underbrace{\begin{bmatrix} h[0] & 0 & 0 & 0 \\ 0 & h[0] & 0 & 0 \\ h[1] & 0 & h[0] & 0 \\ 0 & h[1] & 0 & h[0] \\ 0 & 0 & h[1] & 0 \\ 0 & 0 & 0 & h[1] \end{bmatrix}}_{H(z^2)} \cdot \begin{bmatrix} u[0] \\ u[1] \\ u[2] \\ u[3] \end{bmatrix} \\
 &= \dots = \underbrace{\begin{bmatrix} h[0] & 0 \\ h[1] & h[0] \\ 0 & h[1] \end{bmatrix}}_{H(z)} \cdot \underbrace{\begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 \end{bmatrix}}_{\text{2-fold downsampling}} \cdot \begin{bmatrix} u[0] \\ u[1] \\ u[2] \\ u[3] \end{bmatrix}
 \end{aligned}$$





Identities (II)

- Noble identity (II): (only for rational functions)



Example : N=2

$h[0], h[1], 0, 0, 0, \dots$

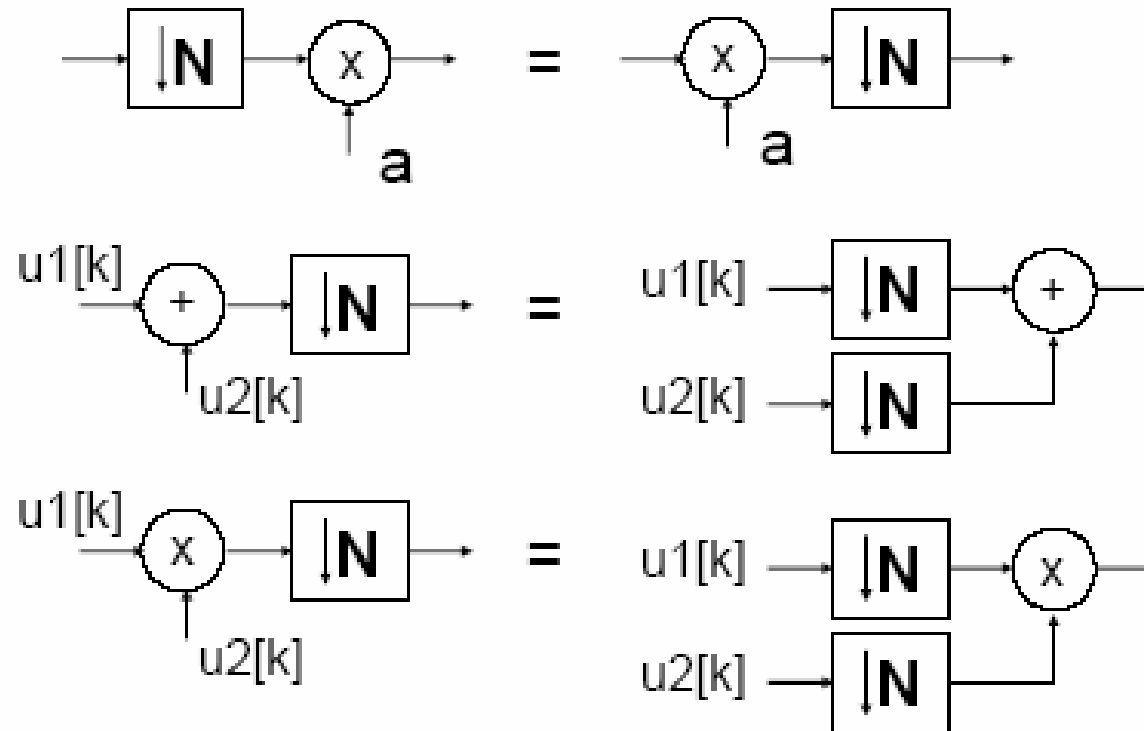
2-fold upsampling

$$\begin{bmatrix} y[0] \\ y[1] \\ y[2] \\ y[3] \\ y[4] \\ y[5] \end{bmatrix} = \begin{bmatrix} 1 & 0 & 0 \\ 0 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 0 \\ 0 & 0 & 1 \\ 0 & 0 & 0 \end{bmatrix} \cdot \begin{bmatrix} h[0] & 0 \\ h[1] & h[0] \\ 0 & h[1] \end{bmatrix} \cdot \begin{bmatrix} u[0] \\ u[1] \end{bmatrix} = \dots = \underbrace{\begin{bmatrix} h[0] & 0 & 0 & 0 \\ 0 & h[0] & 0 & 0 \\ h[1] & 0 & h[0] & 0 \\ 0 & h[1] & 0 & h[0] \\ 0 & 0 & h[1] & 0 \\ 0 & 0 & 0 & h[1] \end{bmatrix}}_{H(z^2)} \cdot \begin{bmatrix} 1 & 0 \\ 0 & 0 \\ 0 & 1 \\ 0 & 0 \end{bmatrix} \cdot \begin{bmatrix} u[0] \\ u[1] \end{bmatrix}$$

2-fold upsampling



Interconnection of Multi-rate Building Blocks

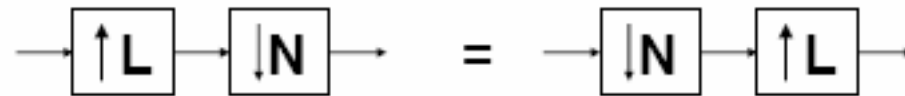


Identities also hold if all decimators are replaced by expanders





Identities (III)



if and only if L and N are *coprime* !!!!!

Example 1: $u[k]=1,2,3,4,5,6,7,8,9,\dots$ (L=2,N=3)

a) 2-fold up: 1,0,2,0,3,0,4,0,... | a) 3-fold down: 1,4,7,...

b) 3-fold down: 1,0,4,0,7,0,... | b) 2-fold up: 1,0,4,0,7,...

Example 2: $u[k]=1,2,3,4,5,6,7,8,9,\dots$ (L=2,N=4)

a) 2-fold up: 1,0,2,0,3,0,4,0,... | a) 4-fold down: 1,5,9,...

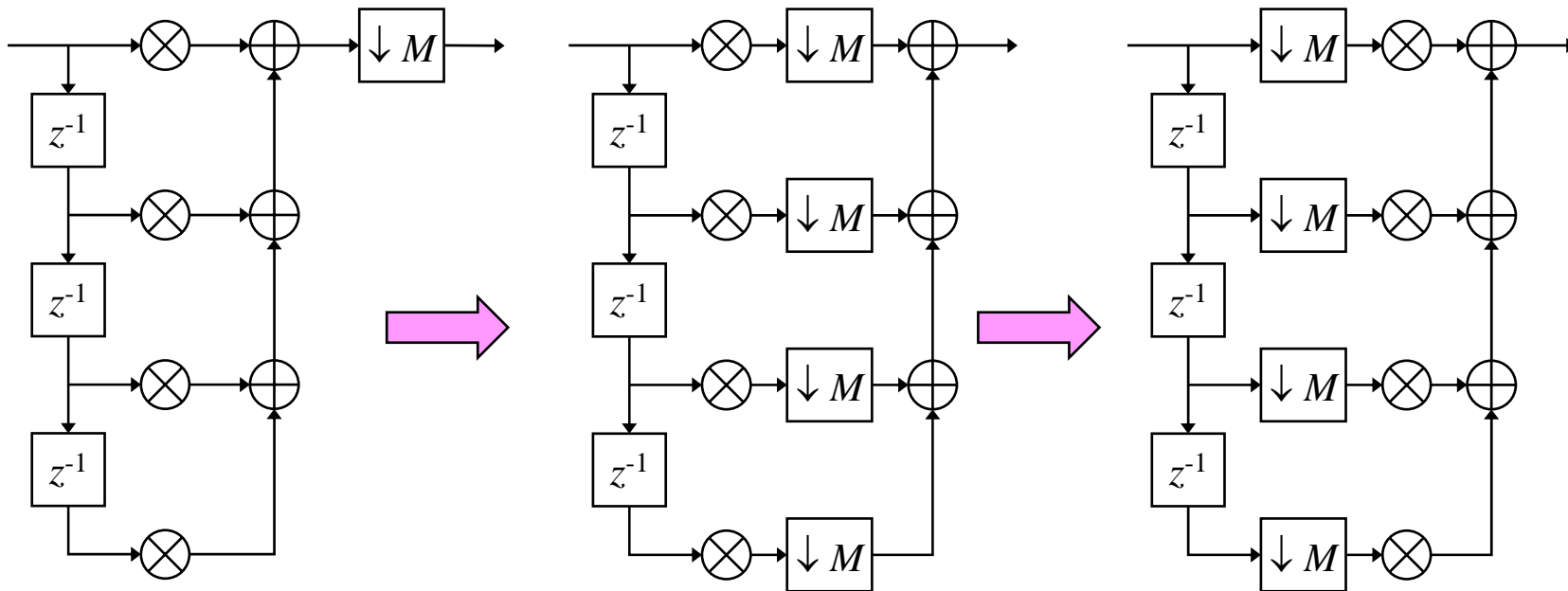
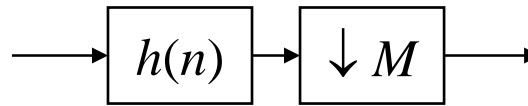
b) 4-fold down: 1,3,5,7,9,... | b) 2-fold up: 1,0,5,0,9,...





Practical Structure

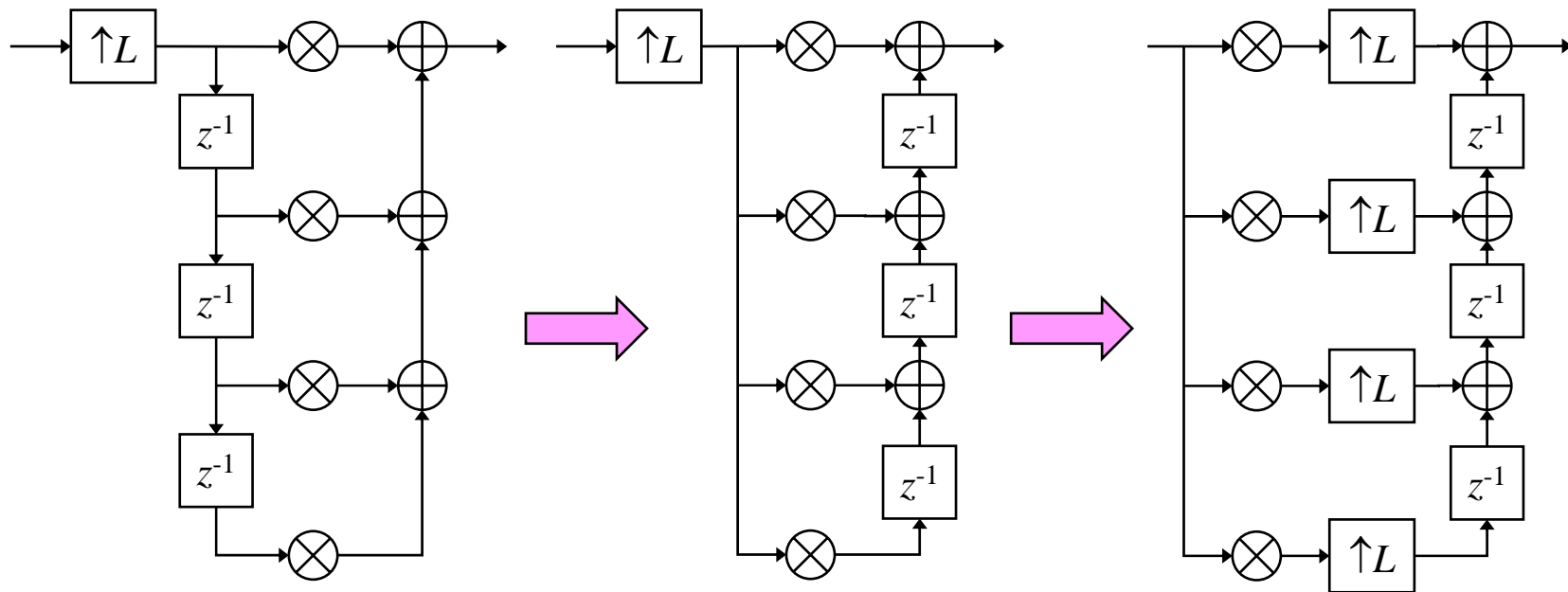
- Decimation





Practical Structure

- Interpolation $\rightarrow \uparrow L \rightarrow h(n) \rightarrow$





An Efficient Decomposition

- Example: 2-fold decomposition

$$\begin{aligned} H(z) &= h[0] + h[1]z^{-1} + h[2]z^{-2} + h[3]z^{-3} + h[4]z^{-4} + h[5]z^{-5} + h[6]z^{-6} \\ &= \underbrace{(h[0] + h[2]z^{-2} + h[4]z^{-4} + h[6]z^{-6})}_{E_0(z^2)} + z^{-1} \underbrace{(h[1] + h[3]z^{-2} + h[5]z^{-4})}_{E_1(z^2)} \end{aligned}$$

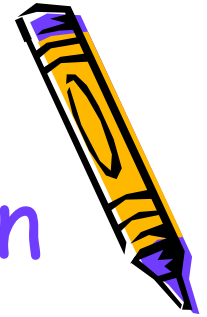
- Example 3-fold decomposition

$$\begin{aligned} H(z) &= h[0] + h[1]z^{-1} + h[2]z^{-2} + h[3]z^{-3} + h[4]z^{-4} + h[5]z^{-5} + h[6]z^{-6} \\ &= \underbrace{(h[0] + h[3]z^{-3} + h[6]z^{-6})}_{E_0(z^3)} + z^{-1} \underbrace{(h[1] + h[4]z^{-3})}_{E_1(z^3)} + z^{-2} \underbrace{(h[2] + h[5]z^{-3})}_{E_2(z^3)} \end{aligned}$$

- General case (N-fold decomposition)

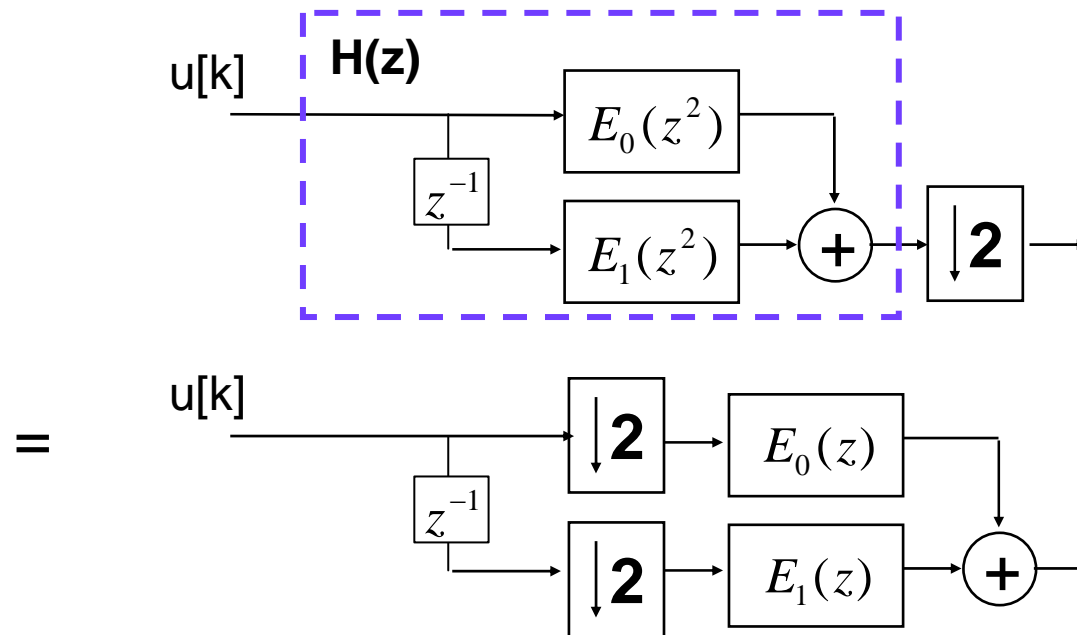
$$H(z) = \sum_{k=-\infty}^{\infty} h[k]z^{-k} = \sum_{l=0}^{N-1} z^{-l} E_l(z^N), \text{ where } E_l(z) = \sum_{k=-\infty}^{\infty} h[Nk + l]z^{-k}$$





Example of Efficient Decomposition

- Efficient implementation of a decimation filter



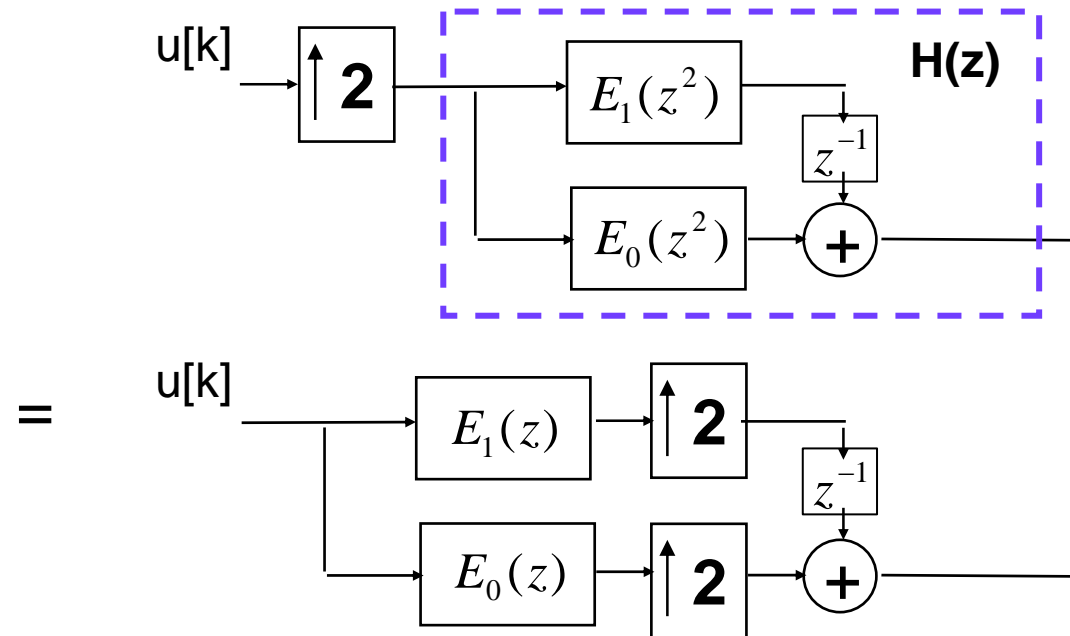
i.e. all filter operations performed at the lowest rate





Example of Efficient Decomposition

- Efficient implementation of an interpolation filter



i.e. all filter operations performed at the lowest rate



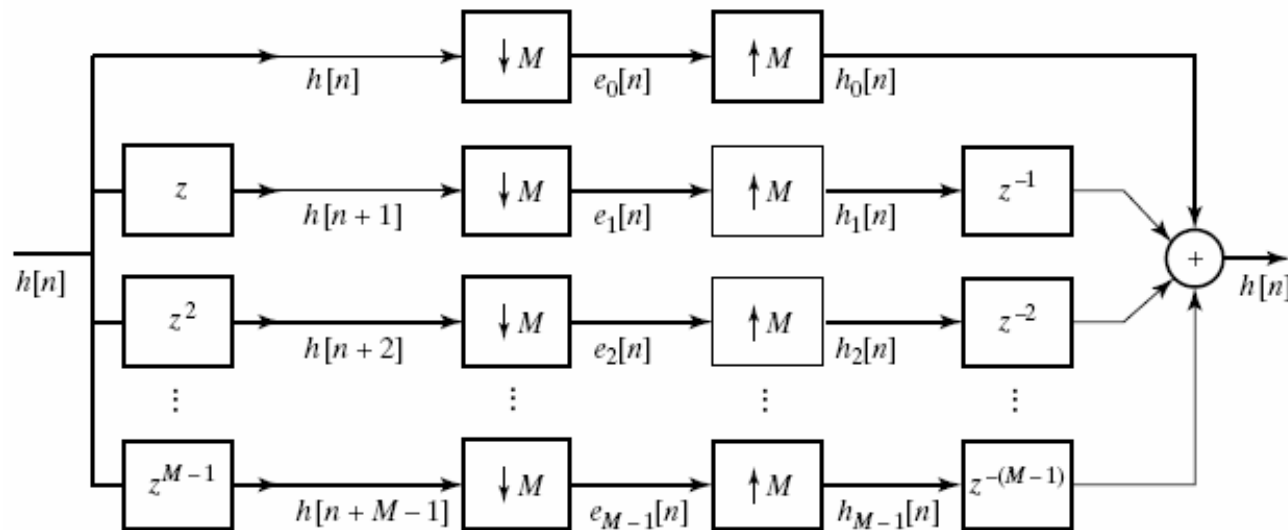


Polyphase Decomposition

- An important advancement in multirate signal processing
 - The polyphase decomposition of a sequence is obtained by representing it as a sequence of M subsequence, each consisting of every M th value of successively delayed versions of the sequence.

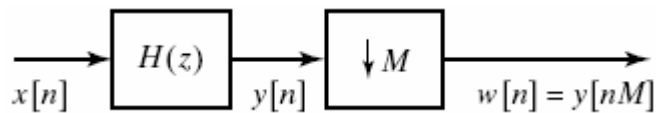
- Example:

$$h[n] = \sum_{k=0}^{M-1} h_k[n-k], \quad \text{with } h_k[n] = \begin{cases} h[n+k], & n = \text{integer multiple of } M \\ 0, & \text{O.W.} \end{cases}$$

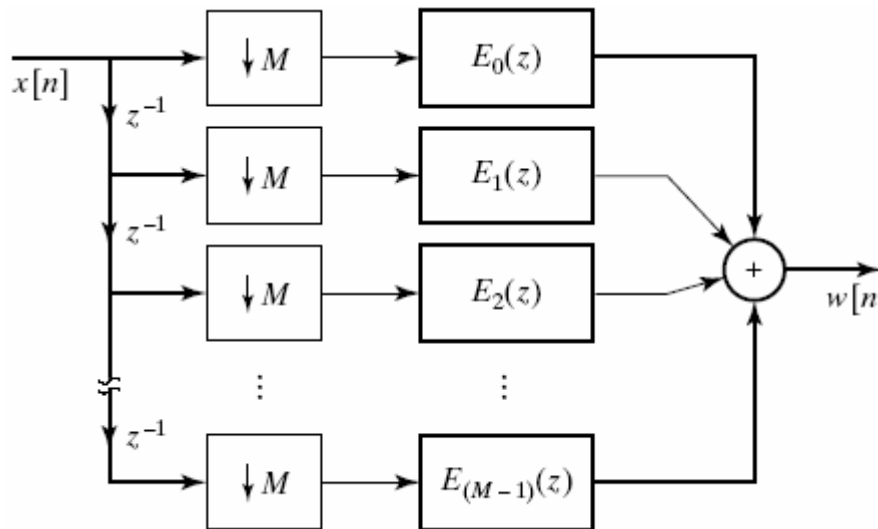




Computation Complexity



N-point FIR



N/M- point FIR

1. N multipliers and N-1 adder
2. One sample in per 1 u.t.
3. One sample out per 1. u.t., but only one of every M outputs is required
4. Computation complexity:
N multiplications, N-1 additions

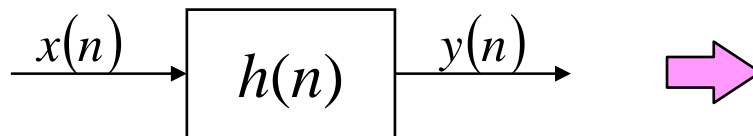
1. For each filter $E_i(z)$, one input is clocked at a rate of $1/M$ u.t.
2. For each filter $E_i(z)$, it requires $(1/M)(N/M)$ multiplications and $(1/M)(N/M-1)$ additions
3. Computation complexity:
N/M multiplications
(N/M-1)+M-1 additions



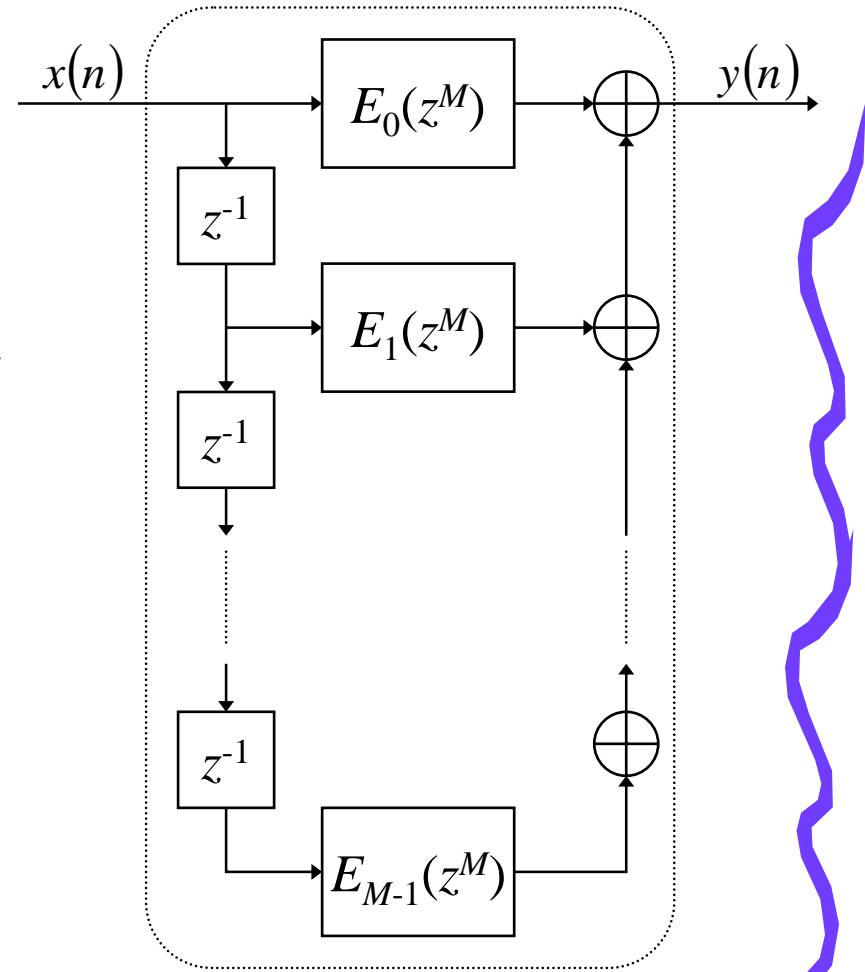


Application: Polyphase FIR Filter

■ Polyphase decomposition

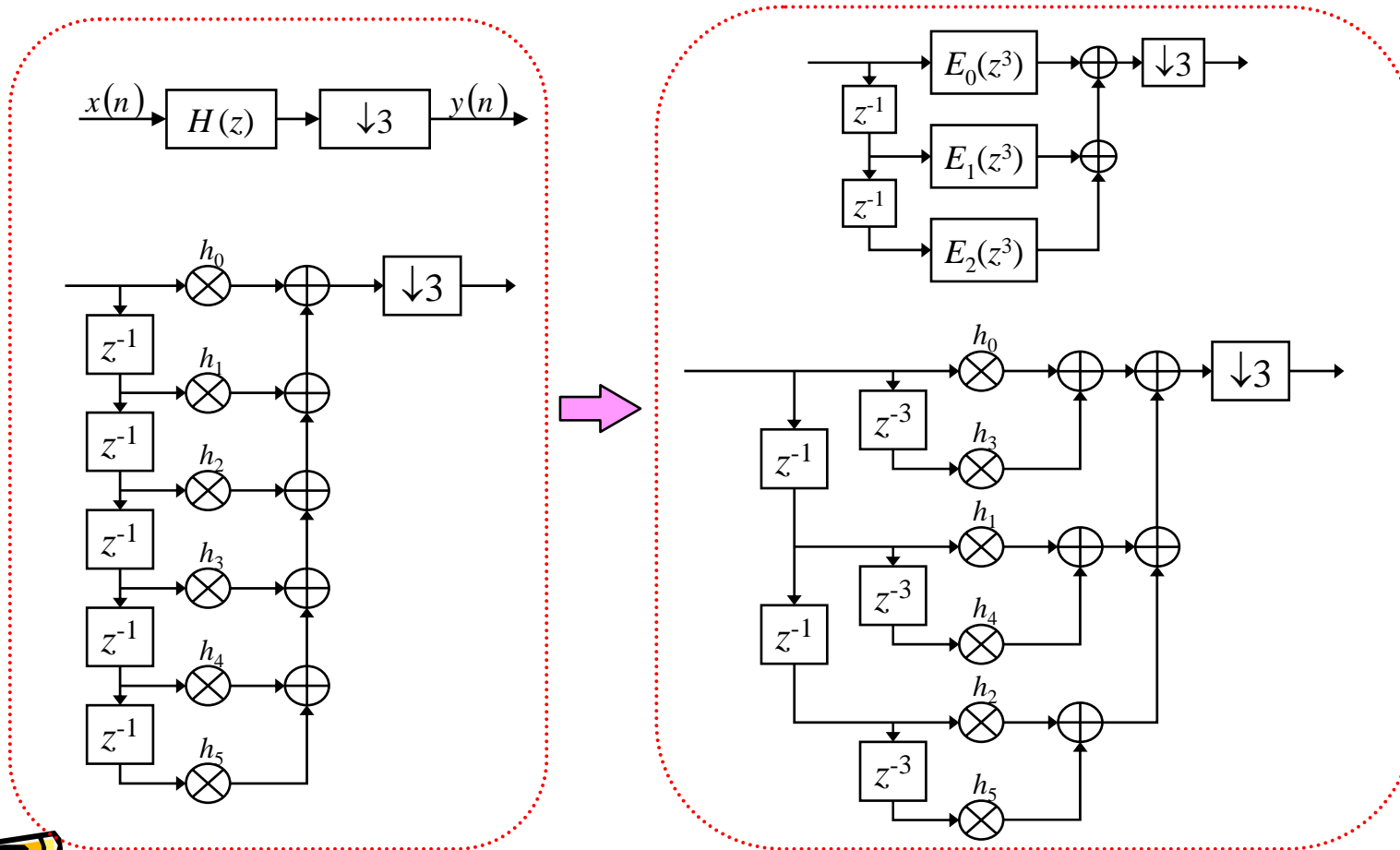


$$H(z) = \sum_{n=-\infty}^{\infty} h(n)z^{-n} \quad \left\{ \begin{array}{l} H(z) = \sum_{l=0}^{M-1} z^{-l} E_l(z^M) \\ E_l(z) = \sum_{n=-\infty}^{\infty} e_l(n)z^{-n} \\ e_l(n) \equiv h(Mn + l) \end{array} \right.$$



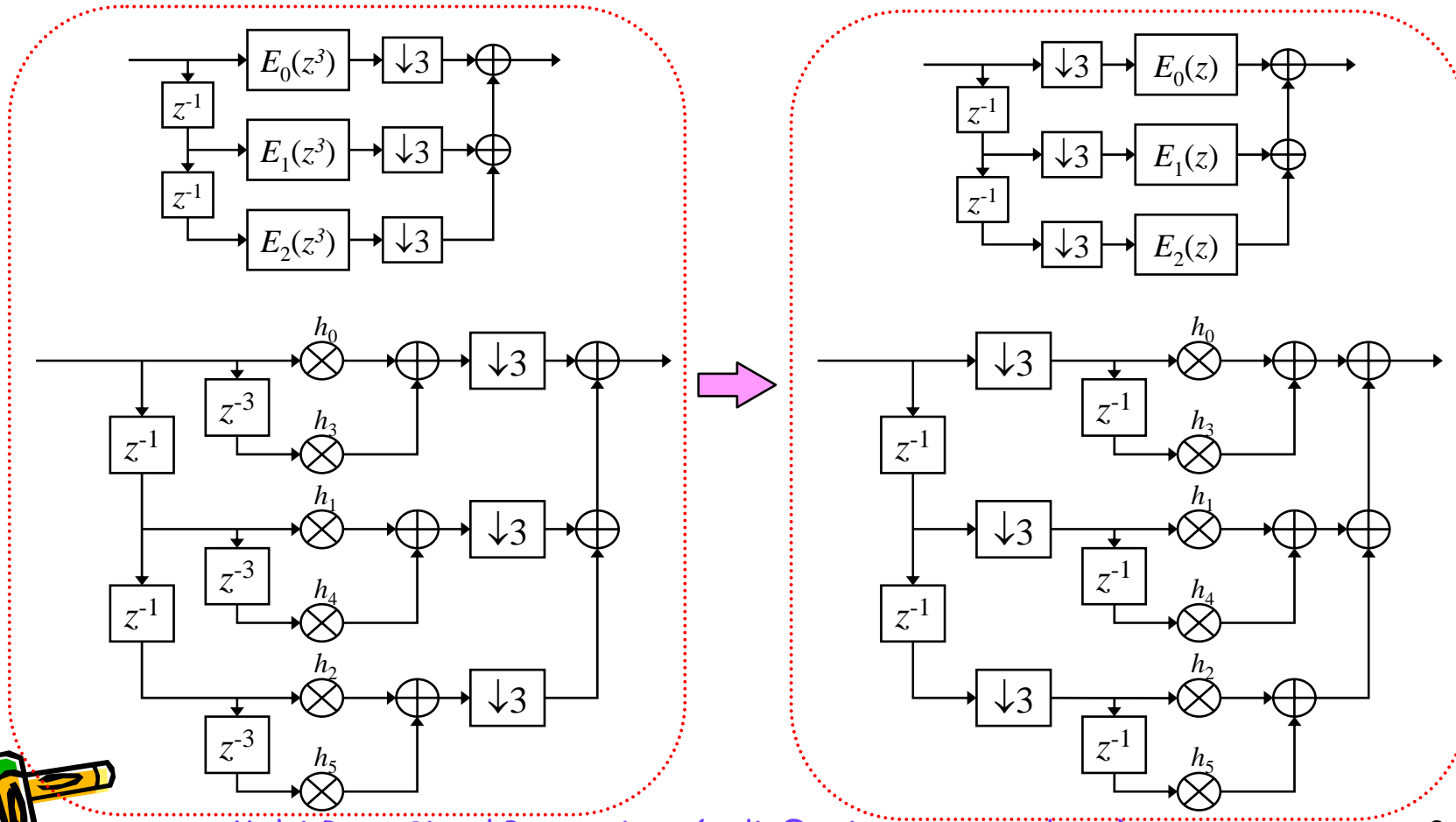


Polyphase FIR Filter



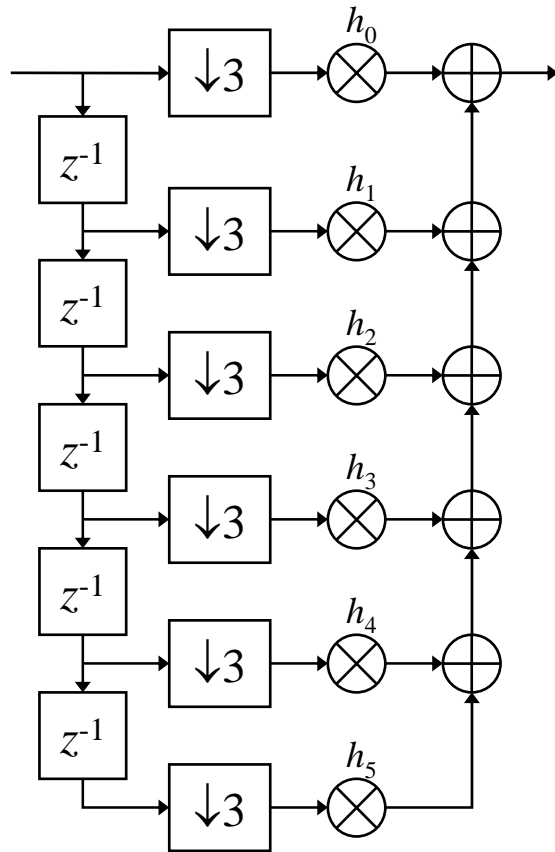


Polyphase FIR Filter

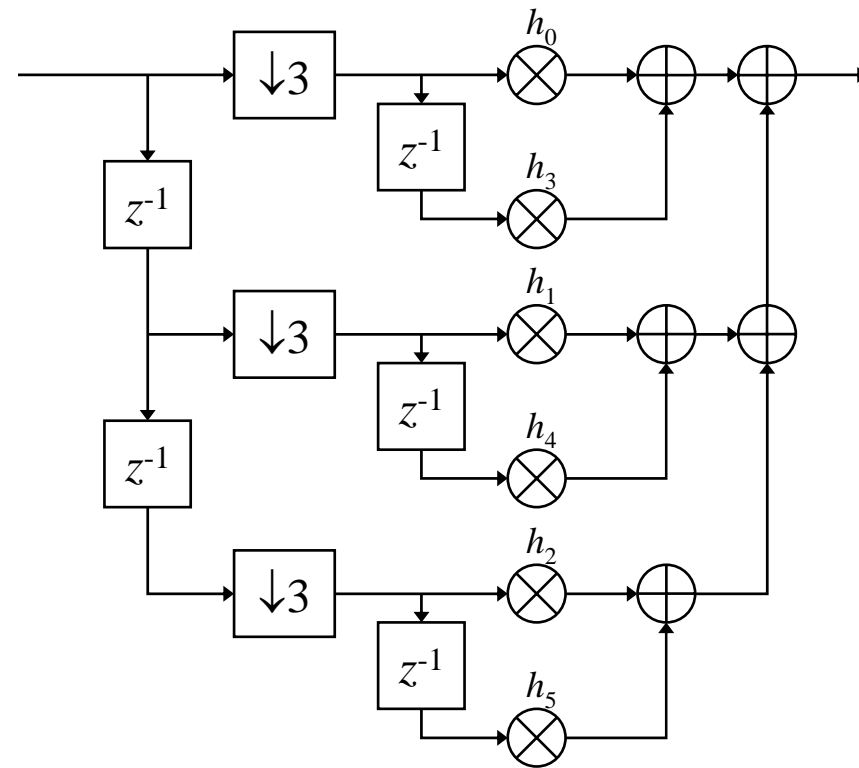




Structure Comparison



Direct implementation



Polyphase implementation





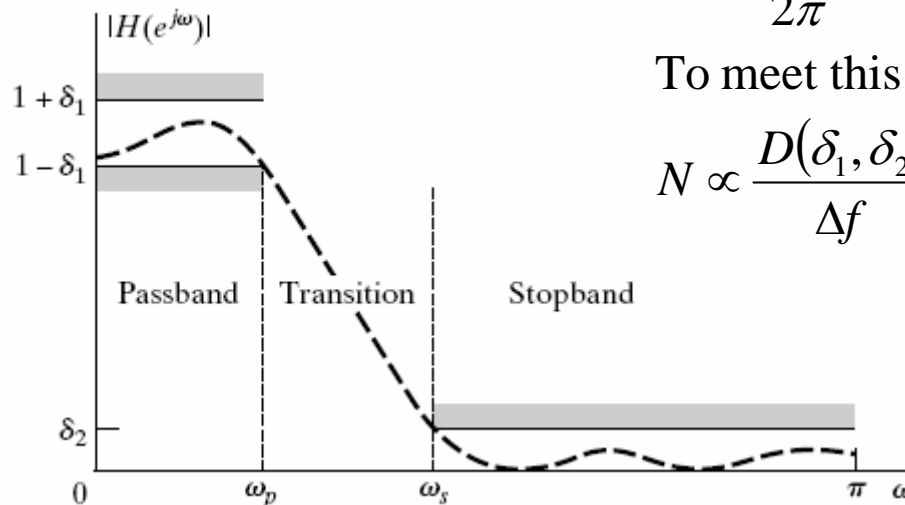
FIR LPF Design Issue

- FIR lowpass filter

$$\Delta f = \frac{\Delta\omega}{2\pi} = \frac{\omega_s - \omega_p}{2\pi} = \text{normalized transition BW}$$

To meet this spec., the FIR filter has order

$$N \propto \frac{D(\delta_1, \delta_2)}{\Delta f}$$



of multiplications $\propto N$
of additions $\propto N$

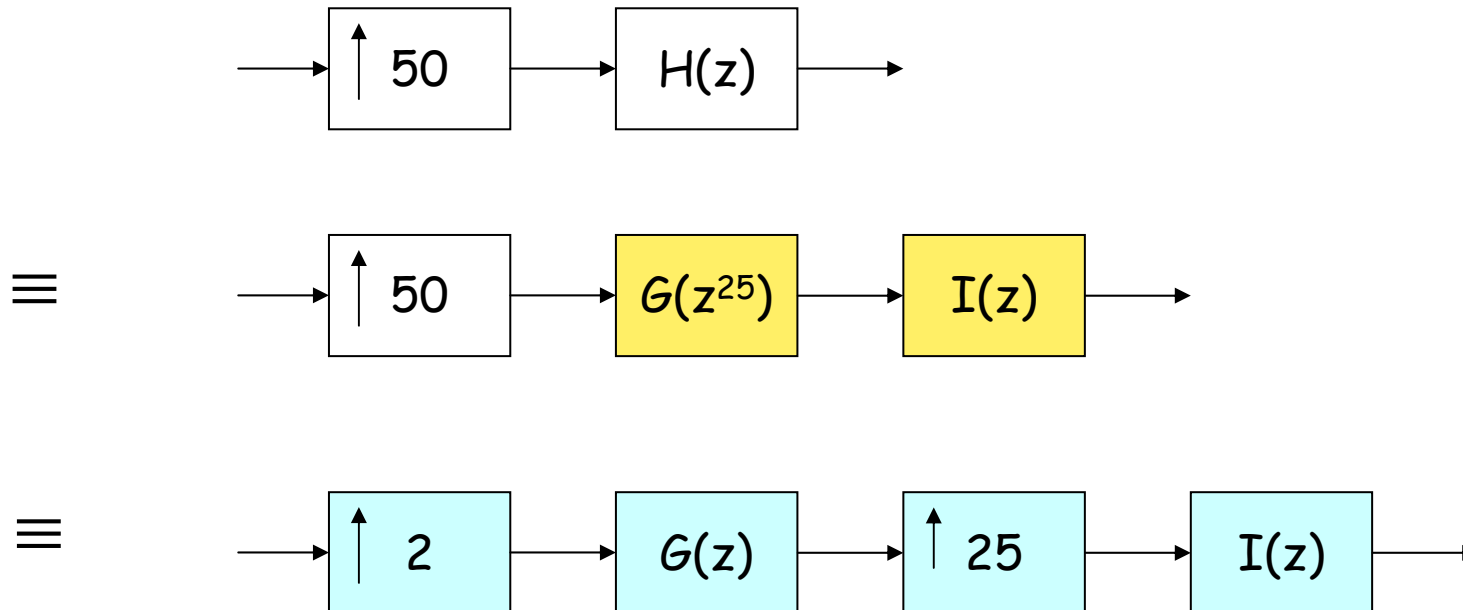
- Large N requires the LPF to have a small stopband edge of π / N





Multistage Implementation

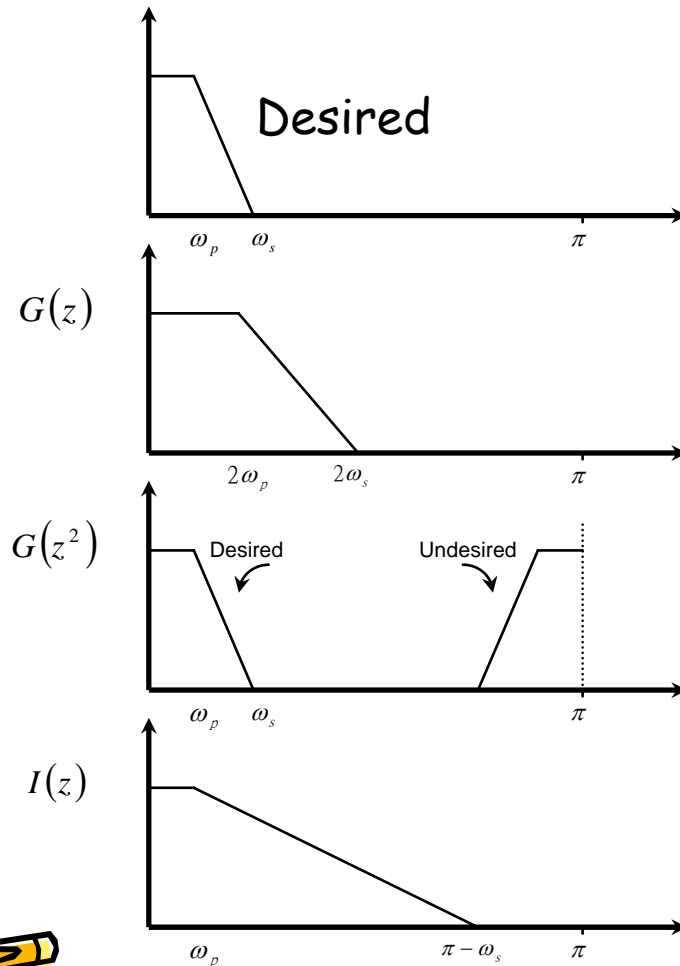
- 2-stage interpolator





Interpolated FIR (IFIR)

with ripple not shown in these figures



Desired narrowband response
Assume required filter order is N .

2-fold Stretched filter
Required filter order is reduced to $N/2$.

Interpolated version of stretched filter
Required filter order is still $N/2$.

Image suppressor
Required filter order is M .
Order $(N/2+M)$ is needed to implement!
 $(N/2+M) \ll N$ for small M



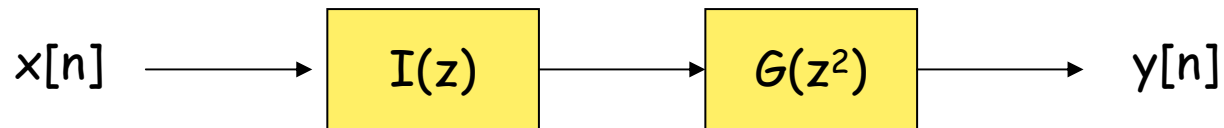


Design of Narrow Band Filter



$\sim N$

(a)



$\sim M$
Small value

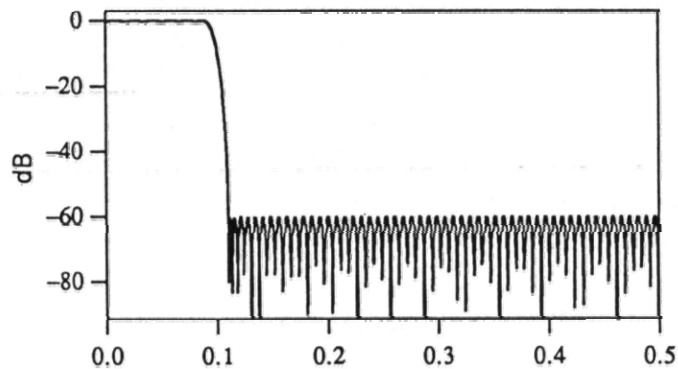
$\sim N/2$

(b)

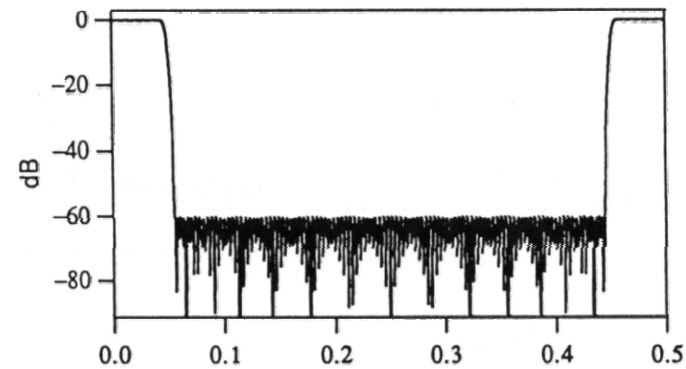




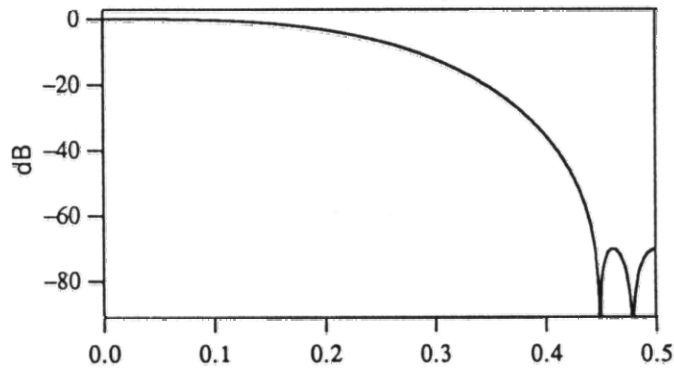
Interpolated FIR (IFIR)



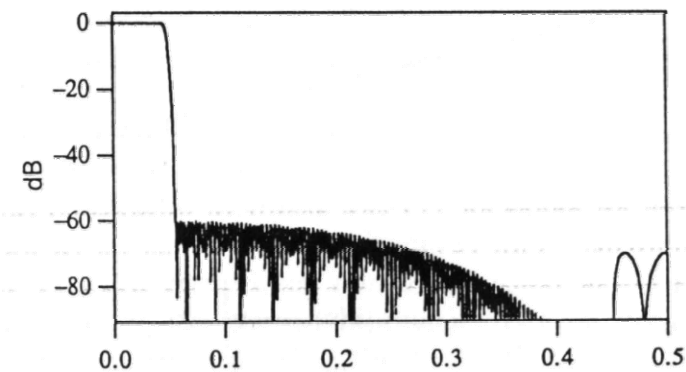
(a) $G(z)$



(a) $G(z^2)$



(b) $I(z)$



(a) $G(z^2)I(z)$





Interpolated FIR (IFIR)

Quantity Compared	Conventional Method	IFIR Method		
		G(z)	I(z)	Total
Filter order	233	131	6	268
Number of Multipliers	<u>117</u>	66	4	<u>70</u>
Number of Adders	233	131	6	137



Remarks

- The overall responses $G(z^2)I(z)$ can have larger ripple sizes. Adjusting ripple sizes is necessary to meet the desired specification.



Questions



- What is the best way to split the integer L (or M) into factors?
- In what order should these factors be arranged?
- **It's hard to answer these questions.** But, it's easy to mention that the multistage implementations result in more efficient systems
- Consider the following FIR LPF design...





Multistage Implementation

- More efficient implementation of the filter (in terms of # computations per time unit)

