

## Solution 2.7

2.7.1

<b>a.</b>	-1391460350
<b>b.</b>	-19629

2.7.2

<b>a.</b>	2903506946
<b>b.</b>	4294947667

2.7.3

<b>a.</b>	AD100002
<b>b.</b>	FFFFFB353

2.7.4

2,7,5

a.	7FFFFFFF
b.	3E8

2.7.6

<b>a.</b>	80000001
<b>b.</b>	FFFFFC18

## **Solution 2.13**

### **2.13.1**

a.	0x57755778
b.	0xFFFFFEDE

### **2.13.2**

a.	0x00005550
b.	0x0000EED0

### **2.13.3**

a.	0x0000AAAA
b.	0x0000BFCD

### **2.13.4**

a.	0x00015B5A
b.	0x000000D0

### **2.13.5**

a.	0xEF000000
b.	0x00000000

### **2.13.6**

a.	0xFFFFFFFF
b.	0x000000F0

## Solution 2.19

### 2.19.1

a.	<pre>compare:     addi \$sp, \$sp, -4     sw   \$ra, 0(\$sp)     add  \$s0, \$a0, \$0     add  \$s1, \$a1, \$0      jal  sub     addi \$t1, \$0, 1     beq  \$v0, \$0, exit     slt  \$t2, \$0, \$v0     bne  \$t2, \$0, exit     addi \$t1, \$0, \$0  exit:     add  \$v0, \$t1, \$0     lw   \$ra, 0(\$sp)     addi \$sp, \$sp, 4     jr   \$ra  sub:     sub  \$v0, \$a0, \$a1     jr   \$ra</pre>
b.	<pre>fib_iter:     addi \$sp, \$sp, -16     sw   \$ra, 12(\$sp)     sw   \$s0, 8(\$sp)     sw   \$s1, 4(\$sp)     sw   \$s2, 0(\$sp)      add  \$s0, \$a0, \$0     add  \$s1, \$a1, \$0     add  \$s2, \$a2, \$0      add  \$v0, \$s1, \$0,     bne  \$s2, \$0, exit      add  \$a0, \$s0, \$s1     add  \$a1, \$s0, \$0     add  \$a2, \$s2, -1     jal  fib_iter  exit:     lw   \$s2, 0(\$sp)     lw   \$s1, 4(\$sp)     lw   \$s0, 8(\$sp)     lw   \$ra, 12(\$sp)     addi \$sp, \$sp, 16     jr   \$ra</pre>

## 2.19.2

a.	compare: addi \$sp, \$sp, -4 sw \$ra, 0(\$sp)  sub \$t0, \$a0, \$a1 addi \$t1, \$0, 1 beq \$t0, \$0, exit slt \$t2, \$0, \$t0 bne \$t2, \$0, exit addi \$t1, \$0, \$0  exit: add \$v0, \$t1, \$0 lw \$ra, 0(\$sp) addi \$sp, \$sp, 4 jr \$ra
b.	Due to the recursive nature of the code, not possible for the compiler to in-line the function call.

## 2.19.3

a.	after calling function compare: old \$sp => 0x7fffffff ??? \$sp => -4 contents of register \$ra  after calling function sub: old \$sp => 0x7fffffff ??? -4 contents of register \$ra \$sp => -8 contents of register \$ra #return to compare
b.	after calling function fib_iter: old \$sp => 0x7fffffff ??? -4 contents of register \$ra -8 contents of register \$s0 -12 contents of register \$s1 \$sp => -16 contents of register \$s2

## 2.19.4

a.	f: addi \$sp,\$sp,-8 sw \$ra,4(\$sp) sw \$s0,0(\$sp) move \$s0,\$a2 jal func move \$a0,\$v0 move \$a1,\$s0 jal func lw \$ra,4(\$sp) lw \$s0,0(\$sp) addi \$sp,\$sp,8 jr \$ra
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<b>b.</b>	<pre> f:    addi    \$sp,\$sp,-12       sw      \$ra,8(\$sp)       sw      \$s1,4(\$sp)       sw      \$s0,0(\$sp)       move   \$s0,\$a1       move   \$s1,\$a2       jal    func       move   \$a0,\$s0       move   \$a1,\$s1       move   \$s0,\$v0       jal    func       add    \$v0,\$v0,\$s0       lw     \$ra,8(\$sp)       lw     \$s1,4(\$sp)       lw     \$s0,0(\$sp)       addi  \$sp,\$sp,12       jr    ra </pre>
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## 2.19.5

<b>a.</b>	We can use the tail-call optimization for the second call to <code>func</code> , but then we must restore <code>\$ra</code> and <code>\$sp</code> before that call. We save only one instruction ( <code>jr \$ra</code> ).
<b>b.</b>	We can NOT use the tail call optimization here, because the value returned from <code>f</code> is not equal to the value returned by the last call to <code>func</code> .

**2.19.6** Register `$ra` is equal to the return address in the caller function, registers `$sp` and `$s3` have the same values they had when function `f` was called, and register `$t5` can have an arbitrary value. For register `$t5`, note that although our function `f` does not modify it, function `func` is allowed to modify it so we cannot assume anything about the value of `$t5` after function `func` has been called.

## Solution 2.25

**2.25.1** Generally, all solutions are similar:

```

lui $t1, top_16_bits
ori $t1, $t1, bottom_16_bits

```

**2.25.2** Jump can go up to 0xFFFFFFF.

<b>a.</b>	no
<b>b.</b>	no

**2.25.3** Range is  $0x604 + 0x1FFFC = 0x0002\ 0600$  to  $0x604 - 0x20000 = 0xFFFFE\ 0604$ .

<b>a.</b>	no
<b>b.</b>	yes

**2.25.4** Range is  $0x0042\ 0600$  to  $0x003E\ 0600$ .

<b>a.</b>	no
<b>b.</b>	no

## 2.25.5

Generally, all solutions are similar:

```
add $t1, $zero, $zero      #clear $t1
addi $t2, $zero, top_8_bits #set top 8b
sll $t2, $t2, 24           #shift left 24 spots
or  $t1, $t1, $t2          #place top 8b into $t1
addi $t2, $zero, nxt1_8_bits #set next 8b
sll $t2, $t2, 16           #shift left 16 spots
or  $t1, $t1, $t2          #place next 8b into $t1
addi $t2, $zero, nxt2_8_bits #set next 8b
sll $t2, $t2, 24           #shift left 8 spots
or  $t1, $t1, $t2          #place next 8b into $t1
ori $t1, $t1, bot_8_bits   #or in bottom 8b
```

## 2.25.6

a.	0x12345678
b.	0x12340000

## 2.25.7

a.	t0 = (0x1234 << 16)   0x5678;
b.	t0 = (t0   0x5678); t0 = 0x1234 << 16;

## Solution 2.39

### 2.39.1

a.	0.86 seconds
b.	0.78 seconds

**2.39.2** Answer is no in all cases. Slows down the computer.

CCT = clock cycle time

ICa = instruction count (arithmetic)

ICls = instruction count (load/store)

ICb = instruction count (branch)

$$\begin{aligned}\text{new CPU time} &= 0.75 \times \text{old ICa} \times \text{CPIa} \times 1.1 \times \text{oldCCT} \\ &\quad + \text{oldICls} \times \text{CPIls} \times 1.1 \times \text{oldCCT} \\ &\quad + \text{oldICb} \times \text{CPIb} \times 1.1 \times \text{oldCCT}\end{aligned}$$

The extra clock cycle time adds sufficiently to the new CPU time such that it is not quicker than the old execution time in all cases.

### 2.39.3

a.	113.16%	121.13%
b.	106.85%	110.64%

### 2.39.4

a.	3
b.	2.65

### 2.39.5

a.	0.6
b.	1.07

### 2.39.6

a.	0.2
b.	0.716666667